

## COMPONENT FUNCTORS

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**Introduction.** Many concrete categories  $\mathcal{Q}$  admit, besides the usual forgetful functor, another such functor  $V: \mathcal{Q} \rightarrow \mathcal{S}$  (where  $\mathcal{S}$  is the category of sets) which may reasonably be termed a "component functor." For example, let  $M$  be a monoid and let  $\mathcal{Q}$  be the category of sets on which  $M$  acts on the left. If  $M$  acts on  $A$ , let  $\sim$  be the equivalence relation generated by  $a \sim ma$  for  $a \in A, m \in M$ . Then the component functor  $V$  takes  $A$  to the set  $A/\sim$  of connected components of  $A$  under the action of  $M$ .

Many such functors are especially interesting in that they have right adjoints  $U: \mathcal{S} \rightarrow \mathcal{Q}$  which are *monadic* (i.e., triplable). In this paper we show that this situation will arise whenever  $\mathcal{Q}$  is *comonadic* over  $\mathcal{S}$ , and exhibit some typical examples.

1. **Statement of the theorem.** Let  $F = (F, \epsilon, \delta)$  be a comonad over  $\mathcal{S}$ . We recall that this means that  $F: \mathcal{S} \rightarrow \mathcal{S}$  is a functor and that  $\epsilon: F \rightarrow 1_{\mathcal{S}}$  and  $\delta: F \rightarrow F^2 = F \circ F$  are natural transformations such that  $F\epsilon \circ \delta = \epsilon F \circ \delta = 1_F$  and  $F\delta \circ \delta = \delta F \circ \delta$ . An algebra over  $F$  is a pair  $(A, \xi)$  consisting of a set  $A$  and a map  $\xi: A \rightarrow FA$  such that  $\epsilon_A \circ \xi = 1_A$  and  $F(\xi) \circ \xi = \delta_A \circ \xi$ . A homomorphism  $f: (A, \xi) \rightarrow (B, \rho)$  is a map  $f: A \rightarrow B$  such that  $\rho \circ f = F(f) \circ \xi$ . Let  $\mathcal{Q}$  be the category of algebras over  $F$ .

We note that  $\mathcal{Q}$  possesses coproducts, which are formed by disjoint union just as in the category of sets. For example, the coproduct of  $(A, \xi)$  and  $(B, \rho)$  is  $(A \amalg B, \omega)$  where the map  $\omega: A \amalg B \rightarrow F(A \amalg B)$  is the composition

$$A \amalg B \xrightarrow{\xi \amalg \rho} F(A) \amalg F(B) \xrightarrow{\gamma} F(A \amalg B)$$

and  $\gamma$  is defined by  $\gamma \circ i_{F(A)} = F(i_A)$ ,  $\gamma \circ i_{F(B)} = F(i_B)$ . If  $X$  is any set, we shall denote by  $X \cdot A$  the coproduct of  $|X|$  copies of  $A$ .

Note also that  $\mathcal{Q}$  has a terminal object  $T$ . Using  $1$  to denote a fixed one-element set, we have  $T = F(1)$  equipped with the map  $\delta_1: T \rightarrow F(T)$ . Now we are ready to state the

**THEOREM.** *The functor  $U: \mathcal{S} \rightarrow \mathcal{Q}$  defined by  $U(X) = X \cdot T$  is monadic.*

The left adjoint  $V$  to  $U$  is called the *component functor* associated with  $\mathcal{Q}$ .

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2. **Proof of the theorem.** First we shall prove that  $U$  has a left adjoint  $V$ , by appealing to the adjoint-functor theorem of Freyd [4]. Observe that a map  $f: (A, \xi) \rightarrow X \cdot T$  is a homomorphism iff the diagram

$$\begin{array}{ccc}
 A & \xrightarrow{f} & X \cdot T \\
 \xi \downarrow & & \downarrow X \cdot \delta_1 \\
 & & X \cdot F(T) \\
 & & \downarrow \gamma \\
 F(A) & \xrightarrow{F(f)} & F(X \cdot T)
 \end{array}$$

commutes. There is a unique homomorphism  $e_A: A \rightarrow T$  which is equal to  $F(\bar{e}) \circ \xi: A \rightarrow F(A) \rightarrow F(1) = T$  where  $\bar{e}: A \rightarrow 1$  is the unique set map. Then for each  $a \in A$ ,  $f(a)$  must be the element  $e_A(a)$  in some summand of  $X \cdot T$  labeled by  $\chi(a) \in X$ . In this way we obtain a function  $\chi: A \rightarrow X$  and by chasing elements in the diagram it is easily seen that the diagram commutes iff for every  $a \in A$ , we have the additional property that

$$F(f)\xi(a) = F(i_{\chi(a)})\delta_1 e_A(a).$$

Using this characterization of homomorphisms  $A \rightarrow X \cdot T$  we shall show that the functor  $U$  preserves limits.

To show that  $U$  preserves equalizers, let  $f, g: X \rightrightarrows Y$  and let  $h: E \rightarrow X$  be  $eq(f, g)$ . We claim that  $h \cdot T: E \cdot T \rightarrow X \cdot T$  is the equalizer in  $\mathcal{A}$  of  $f \cdot T$  and  $g \cdot T$ . For, suppose that  $k: (A, \xi) \rightarrow X \cdot T$  is such that  $(f \cdot T) \circ k = (g \cdot T) \circ k$ . Then there is a unique set map  $j: A \rightarrow E \cdot T$  such that  $k = (h \cdot T) \circ j$ , and we need only show that  $j$  is a homomorphism. Let  $\bar{k}(a)$  be the element  $e_A(a)$  in the summand of  $X \cdot T$  labeled by  $\chi(a)$  and let  $j(a)$  be the element  $e_A(a)$  in the summand of  $E \cdot T$  labeled by  $\eta(a)$ . Then

$$\begin{aligned}
 F(h \cdot T)F(j)\xi(a) &= F(\bar{k})\xi(a) \\
 &= F(i_{\chi(a)})\delta_1 e_A(a) \\
 &= F(h \cdot T)F(i_{\eta(a)})\delta_1 e_A(a).
 \end{aligned}$$

Since  $F(h \cdot T)$  is monic,  $F(j)\xi(a) = F(i_{\eta(a)})\delta_1 e_A(a)$  and thus  $j$  is indeed a homomorphism.

Next, it is rather more difficult to show that  $U$  preserves products; for simplicity we write the proof for binary products but the proof for infinitary products is similar. Thus we must prove that the product in  $\mathcal{A}$  of  $X \cdot T$  and  $Y \cdot T$  is  $(X \times Y) \cdot T$  where  $X \times Y$  is the product

in  $\mathcal{S}$ . Suppose  $f: (A, \xi) \rightarrow X \cdot T$  and  $g: (A, \xi) \rightarrow Y \cdot T$ , where  $f(a)$  is the  $e_A(a)$  in the  $\chi(a)$ -summand and  $g(a)$  is the  $e_A(a)$  in the  $\eta(a)$ -summand. Define  $h: A \rightarrow (X \times Y) \cdot T$  by requiring  $h(a)$  to be  $e_A(a)$  in the summand labeled by  $(\chi(a), \eta(a))$ . Then clearly we need only show that  $h$  is a homomorphism.

First, we note that

$$\begin{aligned} F(f)\xi(a) &= F(i_{\chi(a)})\delta_1 e_A(a) \\ &= F(i_{\chi(a)})\delta_1 F(\bar{e})\xi(a) \\ &= F(i_{\chi(a)})F^2(\bar{e})\delta_A \xi(a) \\ &= F(i_{\chi(a)} \circ F(e) \circ \xi)\xi(a) \\ &= F(i_{\chi(a)} \circ e_A)\xi(a), \end{aligned}$$

and similarly  $F(g)\xi(a) = F(i_{\eta(a)} \circ e_A)\xi(a)$ . To show that  $h$  is a homomorphism, i.e., to show that the diagram

$$\begin{array}{ccc} A & \xrightarrow{h} & (X \times Y) \cdot T \\ & & \downarrow (X \times Y) \cdot \delta_1 \\ \xi \downarrow & & (X \times Y) \cdot F(T) \\ & & \downarrow \gamma \\ F(A) & \xrightarrow{F(h)} & F((X \times Y) \cdot T) \end{array}$$

commutes, we must show that for all  $a \in A$ ,

$$F(h)\xi(a) = F(i_{(\chi(a), \eta(a))} \circ e_A)\xi(a).$$

The result would be immediate if we knew that the map  $(F(p_1), F(p_2)): F((X \times Y) \cdot T) \rightarrow F(X \cdot T) \times F(Y \cdot T)$  is one-to-one. Unfortunately, there are many examples of comonads for which it is *not* one-to-one. However, it will suffice to show that  $(F(p_1), F(p_2))$  is one-to-one *on the image of  $\gamma \circ (X \times Y) \cdot \delta_1$* . In fact, we shall show that the composition

$$k = (F(p_1), F(p_2)) \circ \gamma \circ (X \times Y) \cdot \delta_1: (X \times Y) \cdot T \rightarrow F(X \cdot T) \times F(Y \cdot T)$$

is actually a one-to-one function.

A little computation using the definition of  $\gamma$  shows that  $k = (\gamma_X, \gamma_Y) \circ (X \times Y) \cdot \delta_1$ , where  $\gamma_X: (X \times Y) \cdot F(T) \rightarrow F(X \cdot T)$  is defined by  $\gamma_X \circ i_{(x,y)} = F(i_x)$ ,  $F(T) \rightarrow F(X \cdot T)$  and likewise  $\gamma_Y \circ i_{(x,y)} = F(i_y)$ . Then we have  $k \circ i_{(x,y)} = (F(i_x), F(i_y)) \circ \delta_1: T \rightarrow F(T) \rightarrow F(X \cdot T) \times F(Y \cdot T)$ . But  $\delta_1$ ,  $F(i_x)$ , and  $F(i_y)$  are one-to-one and hence  $k \circ i_{(x,y)}$  is one-to-one for every  $(x, y) \in X \times Y$ .

To show that  $k$  is one-to-one, it now suffices to show that the images of distinct summands under  $k$  are *disjoint*; that is, if  $(x, y) \neq (\bar{x}, \bar{y})$  then the maps  $(F(i_x), F(i_y)) \circ \delta_1$  and  $(F(i_{\bar{x}}), F(i_{\bar{y}})) \circ \delta_1: T \rightarrow F(X \cdot T) \times F(Y \cdot T)$  have disjoint images. In fact, more is true: if  $x \neq \bar{x}$  then the images of  $F(i_x) \circ \delta_1$  and  $F(i_{\bar{x}}) \circ \delta_1$  are disjoint (and similarly if  $y \neq \bar{y}$ ). For, if they were not disjoint, they would continue to fail to be disjoint upon application of the map  $\epsilon_{X \cdot T}$ . But the diagram

$$\begin{array}{ccccc}
 T & \xrightarrow{\delta_1} & F(T) & \xrightleftharpoons[F(i_{\bar{x}})]{F(i_x)} & F(X \cdot T) \\
 & \searrow^{1_T} & \downarrow \epsilon_T & & \downarrow \epsilon_{X \cdot T} \\
 & & T & \xrightleftharpoons[i_{\bar{x}}]{i_x} & X \cdot T
 \end{array}$$

shows that this would be impossible. Hence  $k$  is one-to-one, and we have completed the proof that  $U$  preserves limits.

Now, in order to show that  $U$  has a left adjoint  $V$ , it suffices to show that for each  $(A, \xi)$  there is a family  $\{B_\alpha\}$  of sets such that whenever  $f: (A, \xi) \rightarrow X \cdot T$ , there is a homomorphism  $\eta: (A, \xi) \rightarrow B_\alpha \cdot T$  and a set map  $g: B_\alpha \rightarrow X$  such that  $f = (g \cdot T) \circ \eta$ . But given  $f$ , its image contains at most  $|A|$  elements and thus is contained in  $B \cdot T$  where  $B \subseteq X$  and  $|B| \leq |A|$ . Thus, for a solution set for  $(A, \xi)$  we can take the set of all subsets of  $A$ . This establishes the existence of the component functor  $V: \mathcal{A} \rightarrow \mathcal{S}$ , by Freyd's adjoint functor theorem.

Finally we proceed to show that  $U$  is monadic by appealing to Beck's Triplability Theorem [1], [3]. Let  $d_0, d_1: X \rightrightarrows Y$  and let  $g: Y \rightarrow W$  be  $\text{coeq}(d_0, d_1)$ . Suppose that in  $\mathcal{A}$  we have a diagram

$$\begin{array}{ccccc}
 & & \xrightarrow{d_0 \cdot T} & & \\
 X \cdot T & \xrightleftharpoons[d_1 \cdot T]{} & Y \cdot T & \xrightleftharpoons[s]{} & Z, \\
 & & \xleftarrow{t} & & 
 \end{array}$$

where  $zs = \text{id.}$ ,  $sz = (d_1 \cdot T) \circ t$ ,  $(d_0 \cdot T) \circ t = \text{id.}$ , and  $z \circ (d_0 \cdot T) = z \circ (d_1 \cdot T)$ . Then  $z = \text{coeq}(d_0 \cdot T, d_1 \cdot T)$ , and in this situation Beck's theorem requires that  $Z \cong W \cdot T$ ,  $z \cong g \cdot T$ . These equations are certainly true set-wise, so what we have to show is that the  $F$ -algebra structure on  $Z$  is that of  $W \cdot T$ . But this is easily seen by observing that  $W \cdot T$  does function as the coequalizer of  $d_0 \cdot T$  and  $d_1 \cdot T$  in  $\mathcal{A}$ . This completes the proof of our theorem.

**3. Examples.** A. Let  $M$  be a monoid, let  $T$  be a set on which  $M$  acts, and let  $\text{Mac}(M, T)$  be the category whose objects are sets  $A$  on

which  $M$  acts, equipped with  $M$ -homomorphisms  $e_A: A \rightarrow T$ . A map  $f: (A, e_A) \rightarrow (B, e_B)$  is to be an  $M$ -homomorphism such that  $e_B \circ f = e_A$ .  $\text{Mac}(M, T)$  is comonadic over  $\mathcal{S}$  by [2]. Then  $U(X) = (X \times T, e_{X \times T})$  where  $m(x, t) = (x, mt)$  and  $e_{X \times T} = p_2$ . The component functor takes  $(A, e_A)$  to the set of connected components of  $A$  under the action of  $M$ . (This example was pointed out to the author by John Isbell in a personal communication.)

B. Let  $\mathcal{Q}$  be the category of sets equipped with equivalence relations. Using  $[a]$  to denote the equivalence class of  $a$ , define a map  $f: (A, \sim) \rightarrow (B, \sim)$  to be a function  $f: A \rightarrow B$  such that  $f[a] = [fa]$  for all  $a \in A$ . Then  $\mathcal{Q}$  is comonadic,  $U(X) = (X, =)$ , and  $V(A, \sim) = A/\sim$ .

C. For a more elaborate example, construct a comonad  $(F, \epsilon, \delta)$  as follows. Let  $F(A) = A \times 2^A$ ,  $\epsilon_A = p_1$ , and define  $\delta_A$  by  $\delta_A(a, \alpha) = (a, \alpha, \{(b, \alpha) \mid b \in \alpha\})$ . Then it is easy to compute that an algebra over this monad is essentially a set  $A$  equipped with a function  $\phi_A: A \rightarrow 2^A$  with the property that  $\phi_A(a) = \phi_A(a')$  whenever  $a' \in \phi_A(a)$ . A homomorphism  $f: (A, \phi_A) \rightarrow (B, \phi_B)$  is a function such that  $f(\phi_A(a)) = \phi_B(f(a))$  for all  $a \in A$ . The terminal object is  $T = \{0, 1\}$  where  $\phi_T(0) = \emptyset$  and  $\phi_T(1) = \{1\}$ . The component functor takes  $(A, \phi_A)$  to the set

$$\{a \mid a \in A, \phi_A(a) = \emptyset\} \cup \{\phi_A(a) \mid a \in A, \phi_A(a) \neq \emptyset\}.$$

This case is particularly interesting since it admits another monadic functor  $U: \mathcal{S} \rightarrow \mathcal{Q}$  where  $U(A) = (A \sqcup \{1\}, \phi)$ ,  $\phi(a) = \{a\}$ , and  $\phi(1) = \emptyset$ . The left adjoint to this functor takes  $(A, \phi_A)$  to the set  $\{\phi_A(a) \mid a \in A, \phi_A(a) \neq \emptyset\}$ . Thus, a category can admit reasonable "component functors" other than those for which we have provided a categorical framework.

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